



A9.3 The fetch-decode-execute cycle in detail

Computer architecture · A level · OCR H446 1.1.1, AQA 7517 4.7.3.2,
Eduqas A500QS 2.1 · about 25 min

What this lesson is about

Register transfer notation, the buses at each step, a full register trace, and a stored program that plays notes.

- Register transfer notation
- Fetch
- Decode
- Execute
- A full trace
- Memory-mapped output

Questions

5 marks in all

1. Put the fetch stage in order.

[1 mark]

Number the lines 1 to 4 to put them in the right order.

- ___ MDR $\leftarrow$ [Memory]addressed
- ___ CIR $\leftarrow$ [MDR]
- ___ PC $\leftarrow$ [PC] + 1
- ___ MAR $\leftarrow$ [PC]

Answer:

```
MAR  $\leftarrow$  [PC]
PC  $\leftarrow$  [PC] + 1
MDR  $\leftarrow$  [Memory]addressed
CIR  $\leftarrow$  [MDR]
```

The address goes to the MAR, the PC moves on, memory returns the instruction to the MDR, and it is copied to the CIR.

2. During the fetch, which bus carries the instruction from memory to the MDR?

[1 mark]

- ☐ A Data bus
- ☐ B Address bus
- ☐ C Control bus
- ☐ D None: it is already in the processor

Answer: A. The address travels on the address bus; the instruction comes back on the data bus.

3. What does this program print?

[1 mark]

```
memory = ['LOAD 3', 'ADD 4', 'HALT', 15, 27]
pc = 0
acc = 0
while True:
    mar = pc
    pc = pc + 1
    mdr = memory[mar]
    cir = mdr
    op = cir.split()
    if op[0] == 'HALT':
        break
    mar = int(op[1])
    mdr = memory[mar]
    acc = mdr if op[0] == 'LOAD' else acc + mdr
print(pc, mar, acc)
```

Answer:

3 2 42

HALT is fetched from address 2, leaving the PC at 3 and the MAR at 2; the accumulator holds 15 + 27.

4. The PC holds 12 at the start of a fetch. What does the PC hold while that instruction is executed, if it is not a branch? [1 mark]

- ☐ A 13
- ☐ B 12
- ☐ C 11
- ☐ D It depends on the operand

Answer: A. The PC is incremented during the fetch, so it already points at the next instruction.

5. How does a branch instruction change the flow of a program?

[1 mark]

- ☐ A Its execute step loads a new address into the PC
- ☐ B It changes the MAR
- ☐ C It rewrites the instructions in memory
- ☐ D It sends an interrupt

Answer: A. The next fetch then uses the new address.

The task: fetch in register transfers

The program below already decodes and executes. Write the **fetch** as four register transfers, in the order in the table above, using the variables `pc`, `mar`, `mdr` and `cir`. Straight after the fetch, before decoding, print the four registers in the form `PC=1 MAR=0 MDR=LOAD 7 CIR=LOAD 7`. Memory holds `LOAD 7`, `OUT`, `ADD 8`, `OUT`, `ADD 8`, `OUT`, `HALT`, then the data 300 and 100. `OUT` plays the accumulator for 0.2 seconds, so a working fetch plays 300, 400 and 500 Hz and prints seven lines, the last `PC=7 MAR=6 MDR=HALT CIR=HALT`. The robot does not move.

```
# the two lines every program starts with: the commands, then the robot
from bugbot import *
connect()

memory = ["LOAD 7", "OUT", "ADD 8", "OUT", "ADD 8", "OUT", "HALT", 300, 100]
pc, mar, mdr, cir, acc = 0, 0, 0, "", 0

while True:
    # fetch: four register transfers, then print the registers
    cir = "HALT"

    # decode
    parts = cir.split()
    opcode = parts[0]
    # execute
    if opcode == "LOAD":
        mar = int(parts[1]); mdr = memory[mar]; acc = mdr
    elif opcode == "ADD":
        mar = int(parts[1]); mdr = memory[mar]; acc = acc + mdr
    elif opcode == "OUT":
        tone(acc, 0.2)
    elif opcode == "HALT":
        break
```

The hint students can ask for: Follow the fetch table: where does the address come from, what moves the PC on, which register receives what memory sends back, and where must the instruction go so the MDR is free again. Print only once all four are done.

A solution

```
from bugbot import *
connect()
memory = ["LOAD 7", "OUT", "ADD 8", "OUT", "ADD 8", "OUT", "HALT", 300, 100]
pc, mar, mdr, cir, acc = 0, 0, 0, "", 0

while True:
    mar = pc
    pc = pc + 1
    mdr = memory[mar]
    cir = mdr
    print(f"PC={pc} MAR={mar} MDR={mdr} CIR={cir}")
    parts = cir.split()
    opcode = parts[0]
    if opcode == "LOAD":
        mar = int(parts[1]); mdr = memory[mar]; acc = mdr
    elif opcode == "ADD":
        mar = int(parts[1]); mdr = memory[mar]; acc = acc + mdr
    elif opcode == "OUT":
```

```
    tone(acc, 0.2)
elif opcode == "HALT":
    break
```

Any program that meets the task's checks is marked correct in the simulator; this is one way, not the only way.